

# COMP 122/L Lecture 19

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Slides adapted from Dr. Kyle Dewey

# Overview

- Circuit minimization
  - Boolean algebra
  - Karnaugh maps

# Circuit Minimization

# Motivation

- Unnecessarily large programs: bad
- Unnecessarily large circuits: Very Bad™
  - Why?

# Motivation

- Unnecessarily large programs: bad
- Unnecessarily large circuits: Very Bad™
  - Why?
    - Bigger circuits = bigger chips = higher cost (non-linear too!)
    - Longer circuits = more time needed to move electrons through = slower

# Simplification

- Real-world formulas can often be simplified, according to algebraic rules
  - How might we simplify the following?

$$R = AB + !AB$$

-How might we simplify this?

# Simplification

- Real-world formulas can often be simplified, according to algebraic rules
  - How might we simplify the following?

$$R = AB + !AB$$

$$R = B(A + !A)$$

$$R = B(\text{true})$$

$$R = B$$

# Simplification Trick

- Look for products that differ only in one variable
  - One product has the original variable ( $A$ )
  - The other product has the other variable ( $\bar{A}$ )

$$R = AB + \bar{A}B$$



# Additional Example I

$!ABCD + ABCD + !AB!CD + AB!CD$

# Additional Example 1

$!ABCD + ABCD + !AB!CD + AB!CD$

$BCD(A + !A) + !AB!CD + AB!CD$

# Additional Example I

$!ABCD + ABCD + !AB!CD + AB!CD$

$BCD(A + !A) + !AB!CD + AB!CD$

$BCD + !AB!CD + AB!CD$

# Additional Example I

$$\overline{A}BCD + ABCD + \overline{A}B\overline{C}D + AB\overline{C}D$$
$$BCD(A + \overline{A}) + \overline{A}B\overline{C}D + AB\overline{C}D$$
$$BCD + \overline{A}B\overline{C}D + AB\overline{C}D$$
$$BCD + B\overline{C}D(\overline{A} + A)$$

# Additional Example I

$$\overline{A}BCD + ABCD + \overline{A}B\overline{C}D + AB\overline{C}D$$

$$BCD(A + \overline{A}) + \overline{A}B\overline{C}D + AB\overline{C}D$$

$$BCD + \overline{A}B\overline{C}D + AB\overline{C}D$$

$$BCD + B\overline{C}D(\overline{A} + A)$$

$$BCD + B\overline{C}D$$

# Additional Example 1

$$\overline{A}BCD + ABCD + \overline{A}B\overline{C}D + AB\overline{C}D$$
$$BCD(A + \overline{A}) + \overline{A}B\overline{C}D + AB\overline{C}D$$
$$BCD + \overline{A}B\overline{C}D + AB\overline{C}D$$
$$BCD + B\overline{C}D(\overline{A} + A)$$
$$BCD + B\overline{C}D$$
$$BD(C + \overline{C})$$

# Additional Example 1

$$!ABCD + ABCD + !AB!CD + AB!CD$$
$$BCD(A + !A) + !AB!CD + AB!CD$$
$$BCD + !AB!CD + AB!CD$$
$$BCD + B!CD(!A + A)$$
$$BCD + B!CD$$
$$BD(C + !C)$$

BD

# Additional Example 2

$!A!BC + A!B!C + !ABC + !AB!C + A!BC$



# Additional Example 2

$!A!BC + A!B!C + !ABC + !AB!C + A!BC$

$!A!BC + A!BC + A!B!C + !ABC + !AB!C$

# Additional Example 2

$!A!BC + A!B!C + !ABC + !AB!C + A!BC$

$!A!BC + A!BC + A!B!C + !ABC + !AB!C$

$!BC(A + !A) + A!B!C + !ABC + !AB!C$

# Additional Example 2

$!A!BC + A!B!C + !ABC + !AB!C + A!BC$

$!A!BC + A!BC + A!B!C + !ABC + !AB!C$

$!BC(A + !A) + A!B!C + !ABC + !AB!C$

$!BC + A!B!C + !ABC + !AB!C$

# Additional Example 2

$!A!BC + A!B!C + !ABC + !AB!C + A!BC$

$!A!BC + A!BC + A!B!C + !ABC + !AB!C$

$!BC(A + !A) + A!B!C + !ABC + !AB!C$

$!BC + A!B!C + !ABC + !AB!C$

$!BC + A!B!C + !AB(C + !C)$

# Additional Example 2

$!A!BC + A!B!C + !ABC + !AB!C + A!BC$

$!A!BC + A!BC + A!B!C + !ABC + !AB!C$

$!BC(A + !A) + A!B!C + !ABC + !AB!C$

$!BC + A!B!C + !ABC + !AB!C$

$!BC + A!B!C + !AB(C + !C)$

$!BC + A!B!C + !AB$

# De Morgan's Laws

Also potentially useful for simplification

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$$\neg (A + B) = \neg A \neg B$$

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Also potentially useful for simplification

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$$\neg (A + B) = \neg A \neg B$$

$$\neg (AB) = \neg A + \neg B$$



# De Morgan Example

$$\neg (X + Y) \neg (!X + Z)$$

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$$\neg (X + Y) \neg (!X + Z)$$

$\neg A$

$\neg B$

Has the overall form of  $\neg A \neg B$

# De Morgan Example

$$\neg (X + Y) \neg (\neg X + Z)$$

$\neg A$

$\neg B$

Has the overall form of  $\neg A \neg B$

# De Morgan Example

$$\neg (X + Y) \neg (!X + Z)$$

$\neg A$

$\neg B$

From De Morgan's Law:

$$\neg (A + B) = \neg A \neg B$$

Has the overall form of  $\neg A \neg B$

# De Morgan Example

$$\neg (X + Y) \neg (!X + Z)$$

$$\neg A \quad \neg B$$

From De Morgan's Law:

$$\neg (A + B) = \neg A \neg B$$

$$\neg (X + Y + !X + Z)$$

Has the overall form of  $\neg A \neg B$

# De Morgan Example

$$\neg (X + Y) \neg (!X + Z)$$

$\neg A$

$\neg B$

From De Morgan's Law:

$$\neg (A + B) = \neg A \neg B$$

$$\neg (X + Y + !X + Z)$$

$$\neg (X + \color{red}{!X} + \color{red}{Y} + Z)$$

Has the overall form of  $\neg A \neg B$

# De Morgan Example

$$\neg (X + Y) \neg (!X + Z)$$

$\neg A$

$\neg B$

From De Morgan's Law:

$$\neg (A + B) = \neg A \neg B$$

$$\neg (X + Y + !X + Z)$$

$$\neg (X + !X + Y + Z)$$

$$\neg (\text{true} + Y + Z)$$

Has the overall form of  $\neg A \neg B$

# De Morgan Example

$$\neg (X + Y) \neg (!X + Z)$$

$\neg A$

$\neg B$

**From De Morgan's Law:**

$$\neg (A + B) = \neg A \neg B$$

$$\neg (X + Y + !X + Z)$$

$$\neg (X + !X + Y + Z)$$

$$\neg (\text{true} + Y + Z)$$

$$\neg (\text{true})$$

Has the overall form of  $\neg A \neg B$



# De Morgan Example

$$\neg (X + Y) \neg (!X + Z)$$
$$\neg A \qquad \qquad \neg B$$

**From De Morgan's Law:**

$$\neg (A + B) = \neg A \neg B$$
$$\neg (X + Y + !X + Z)$$
$$\neg (X + !X + Y + Z)$$
$$\neg (\text{true} + Y + Z)$$
$$\neg (\text{true})$$

false

# Boolean Operators (Truth Table)

AND $X Y$			OR $X + Y$			Complement (opposite) $\bar{X}$	
X	Y	X AND Y	X	Y	X OR Y	X	NOT X
0	0	0	0	0	0	0	1
0	1	0	0	1	1	1	0
1	0	0	1	0	1		
1	1	1	1	1	1		

# Boolean Algebra

Name	AND form	OR form
Identity law	$1A = A$	$0 + A = A$
Null law	$0A = 0$	$1 + A = 1$
Idempotent law	$AA = A$	$A + A = A$
Inverse law	$A\bar{A} = 0$	$A + \bar{A} = 1$
Commutative law	$AB = BA$	$A + B = B + A$
Associative law	$(AB)C = A(BC)$	$(A + B) + C = A + (B + C)$
Distributive law	$A + BC = (A + B)(A + C)$	$A(B + C) = AB + AC$
Absorption law	$A(A + B) = A$	$A + AB = A$
De Morgan's law	$\overline{AB} = \bar{A} + \bar{B}$	$\overline{A + B} = \bar{A}\bar{B}$

# Scaling Up

- Performing this sort of algebraic manipulation by hand can be tricky
- We can use *Karnaugh maps* to make it immediately apparent as to what can be simplified

**Karnaugh Maps - Rules of Simplification:**

<http://www.ee.surrey.ac.uk/Projects/Labview/minimisation/karrules.html>

# Example

$$R = AB + !AB$$

-Start with the sum of products

# Example

$$R = AB + !AB$$

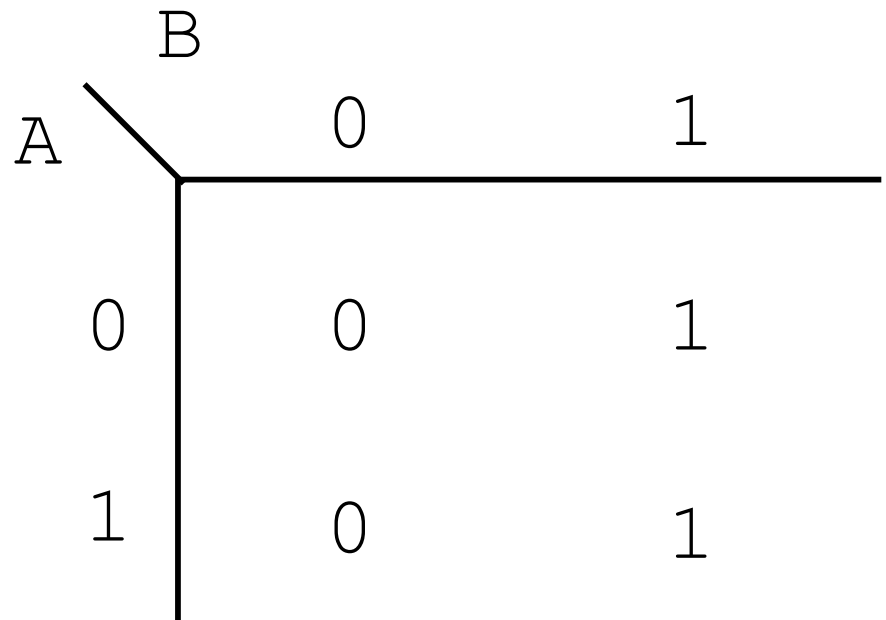
A	B	O
0	0	0
0	1	1
1	0	0
1	1	1

-Build the truth table

# Example

$$R = AB + !AB$$

A	B	O
0	0	0
0	1	1
1	0	0
1	1	1

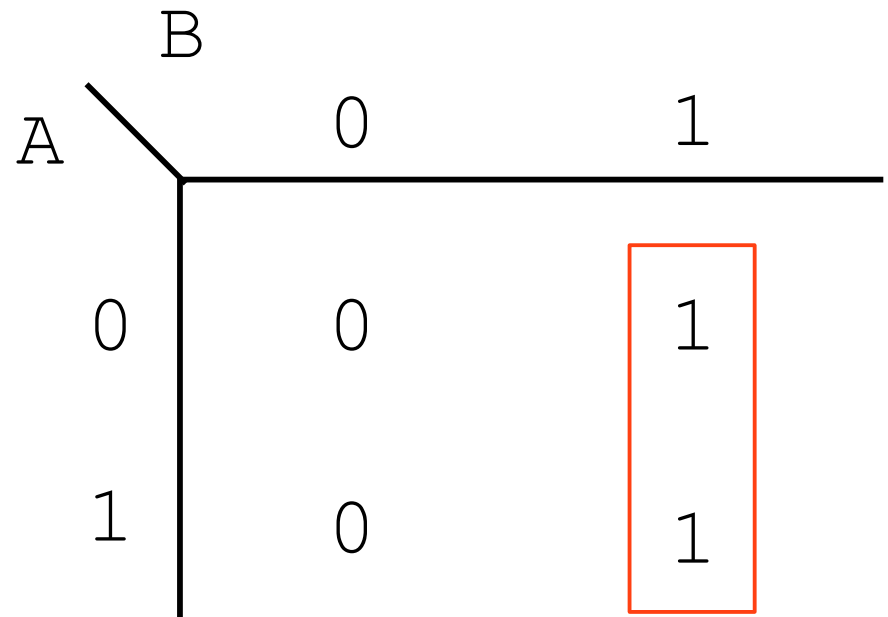


-Build the K-map

# Example

$$R = AB + \neg AB$$

A	B	O
0	0	0
0	1	1
1	0	0
1	1	1



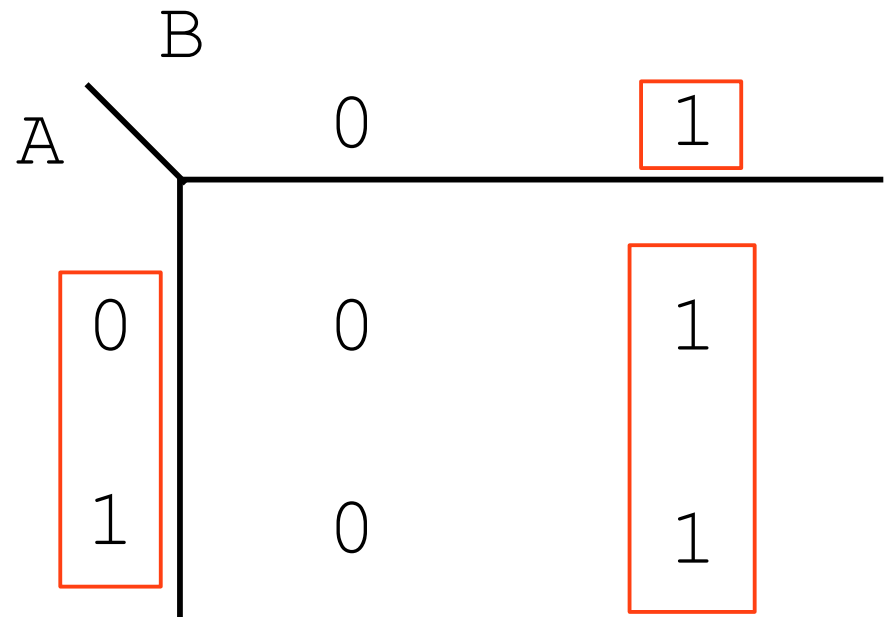
-Group adjacent (row or column-wise, NOT diagonal) 1's in powers of two (groups of 2, 4, 8...)



# Example

$$R = AB + !AB$$

A	B	O
0	0	0
0	1	1
1	0	0
1	1	1

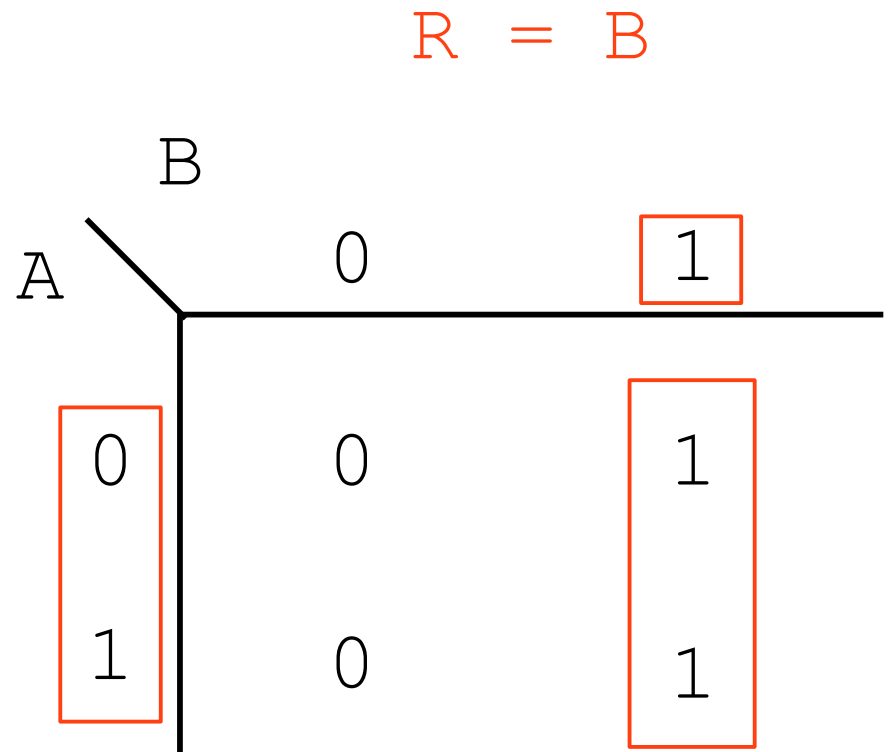


- The values that stay the same are saved, the rest are discarded
- This works because this means that the inputs that differ are irrelevant to the final value, and so they can be removed

# Example

$$R = AB + !AB$$

A	B	O
0	0	0
0	1	1
1	0	0
1	1	1



- The values that stay the same are saved, the rest are discarded
- This works because this means that the inputs that differ are irrelevant to the final value, and so they can be removed

# Three Variables

- We can scale this up to three variables, by combining two variables on one axis
- The combined axis must be arranged such that only one bit changes per position

A	BC			
	00	01	11	10
0	?	?	?	?
1	?	?	?	?

# Three Variable Example

$$R = !A!BC + !ABC + A!BC + ABC$$

-Start with this formula

$$R = \neg A \neg B C + \neg A B C + A \neg B C + A B C$$

A	B	C	R
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

-Build the truth table

$$R = \bar{A}\bar{B}C + \bar{A}BC + A\bar{B}C + ABC$$

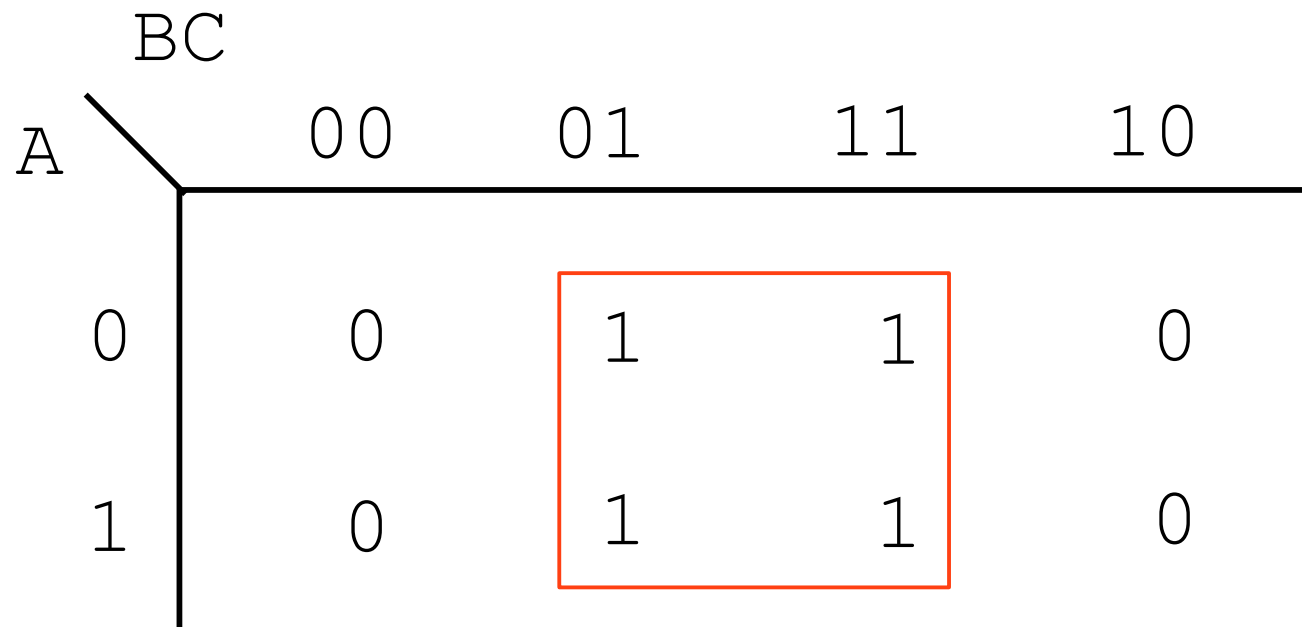
A	B	C	R
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

		BC			
		00	01	11	10
A	0	0	1	1	0
	1	0	1	1	0

-Build the K-map

$$R = !A!BC + !ABC + A!BC + ABC$$

A	B	C	R
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1



-Select the biggest group possible, in this case a square

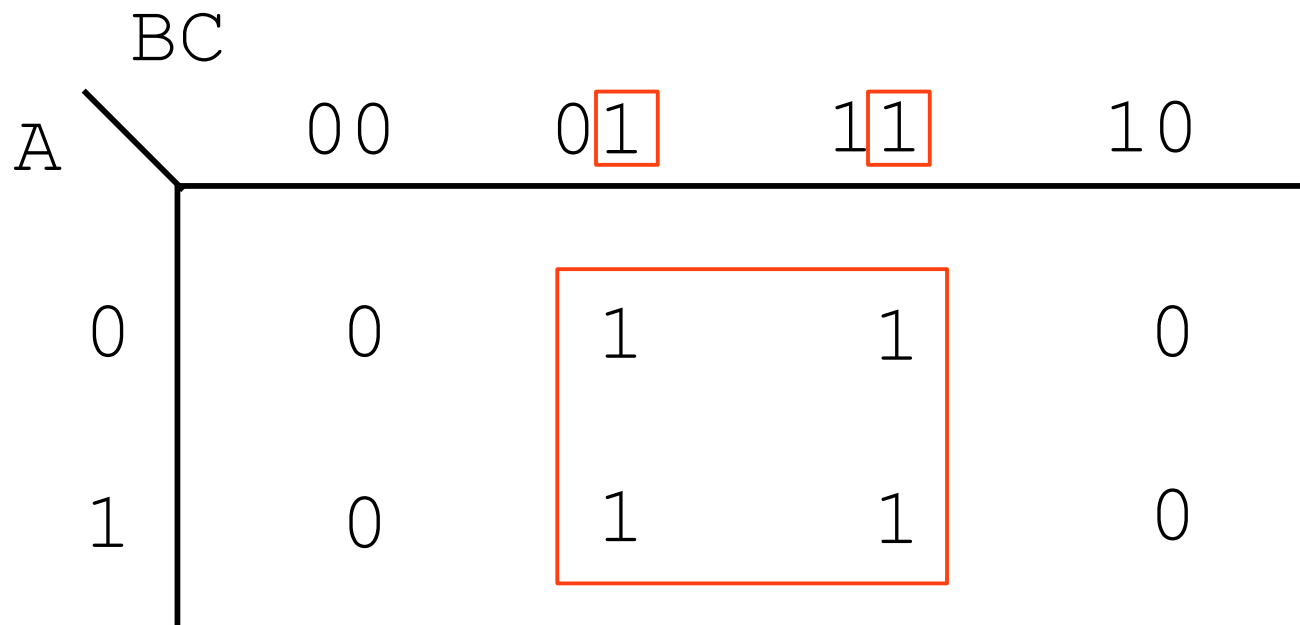
-In order to get the most minimal circuit, we must always select the biggest groups possible



$$R = \neg A \neg B C + \neg A B C + A \neg B C + A B C$$

A	B	C	R
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

$$R = C$$



-Save the ones that stay the same in a group, discarding the rest

# Another Three Variable Example

$$R = !A!B!C + !A!BC + !ABC + \\ !AB!C + A!B!C + AB!C$$

-Start with this formula

$$R = \neg A \neg B \neg C + \neg A \neg B C + \neg A B \neg C + \neg A B C + A \neg B \neg C + A \neg B C + A B \neg C$$

A	B	C	R
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	0

-Build the truth table

$$R = \overline{A}\overline{B}\overline{C} + \overline{A}\overline{B}C + \overline{A}BC + \overline{A}B\overline{C} + A\overline{B}\overline{C} + AB\overline{C}$$

A	B	C	R
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	0

		BC			
		00	01	11	10
A	0	1	1	1	1
	1	1	0	0	1

-Build the K-map

$$R = \overline{A}\overline{B}\overline{C} + \overline{A}\overline{B}C + \overline{A}BC + \overline{A}BC + \overline{A}B\overline{C} + A\overline{B}\overline{C} + AB\overline{C}$$

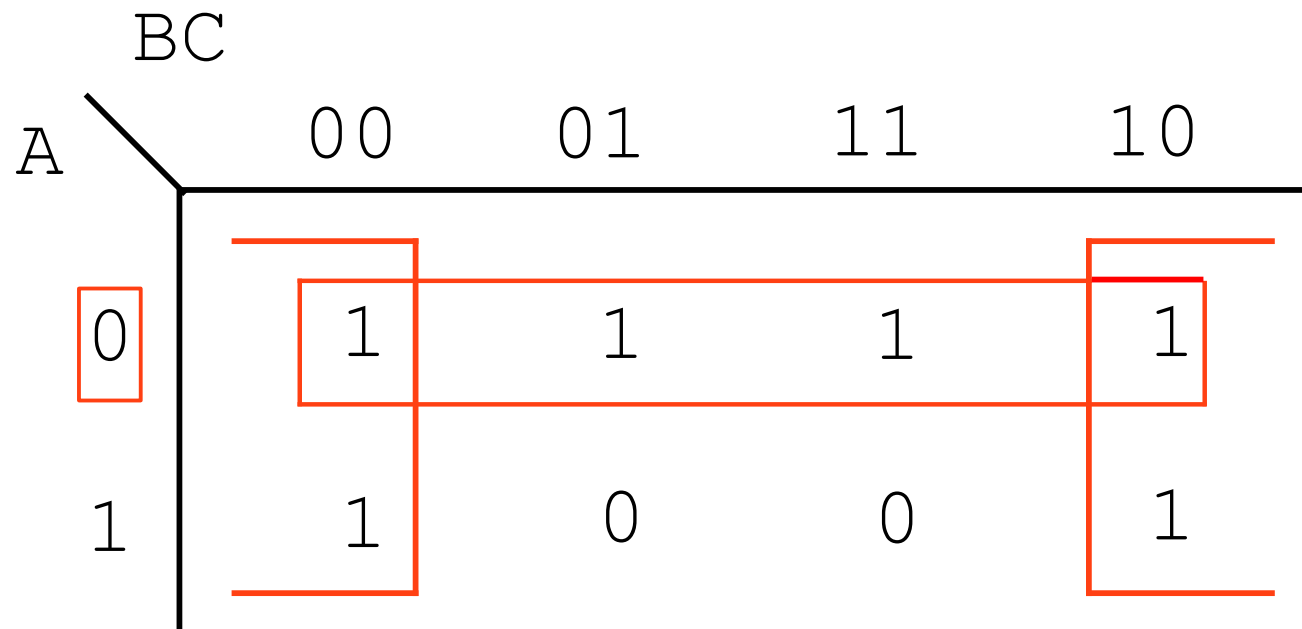
A	B	C	R
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	0

		BC			
		00	01	11	10
A	0	1	1	1	1
	1	1	0	0	1

- Select the biggest groups possible
- Note that the values “wrap around” the table

$$R = \overline{A}\overline{B}\overline{C} + \overline{A}\overline{B}C + \overline{A}BC + \overline{A}BC + A\overline{B}\overline{C} + AB\overline{C}$$

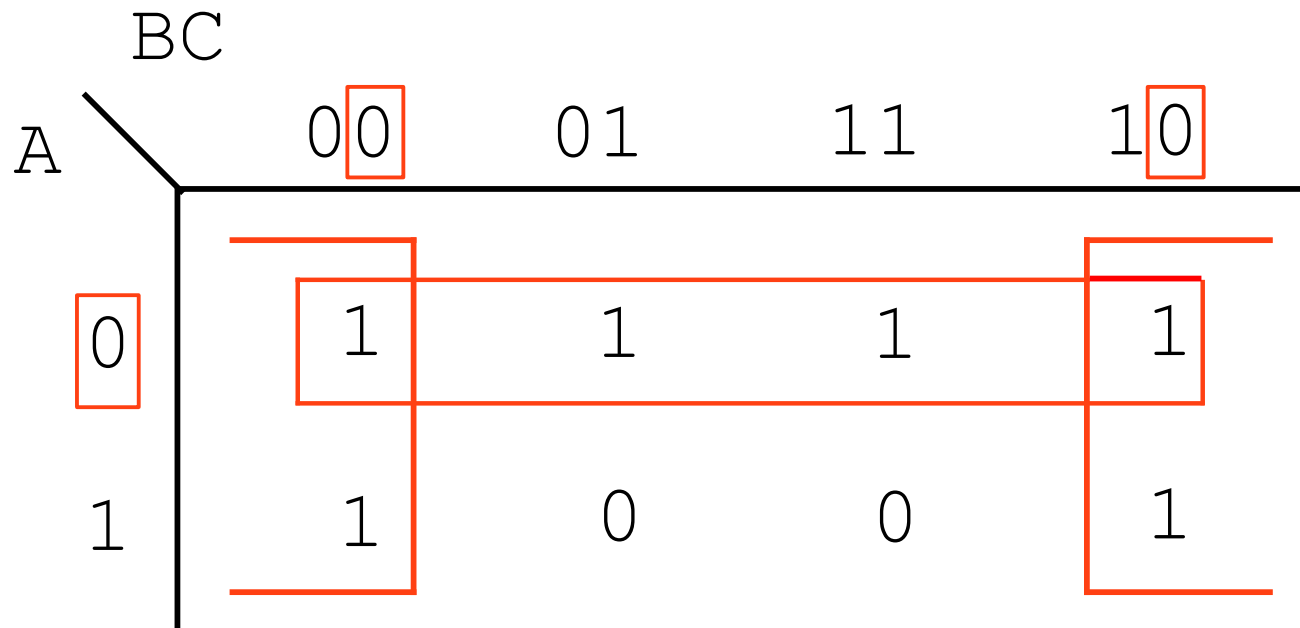
A	B	C	R
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	0



- Save the ones that stay the same in a group, discarding the rest
- This must be done for each group

$$R = \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}C + \bar{A}BC + \bar{A}BC + \bar{A}B\bar{C} + A\bar{B}\bar{C} + AB\bar{C}$$

A	B	C	R
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	0



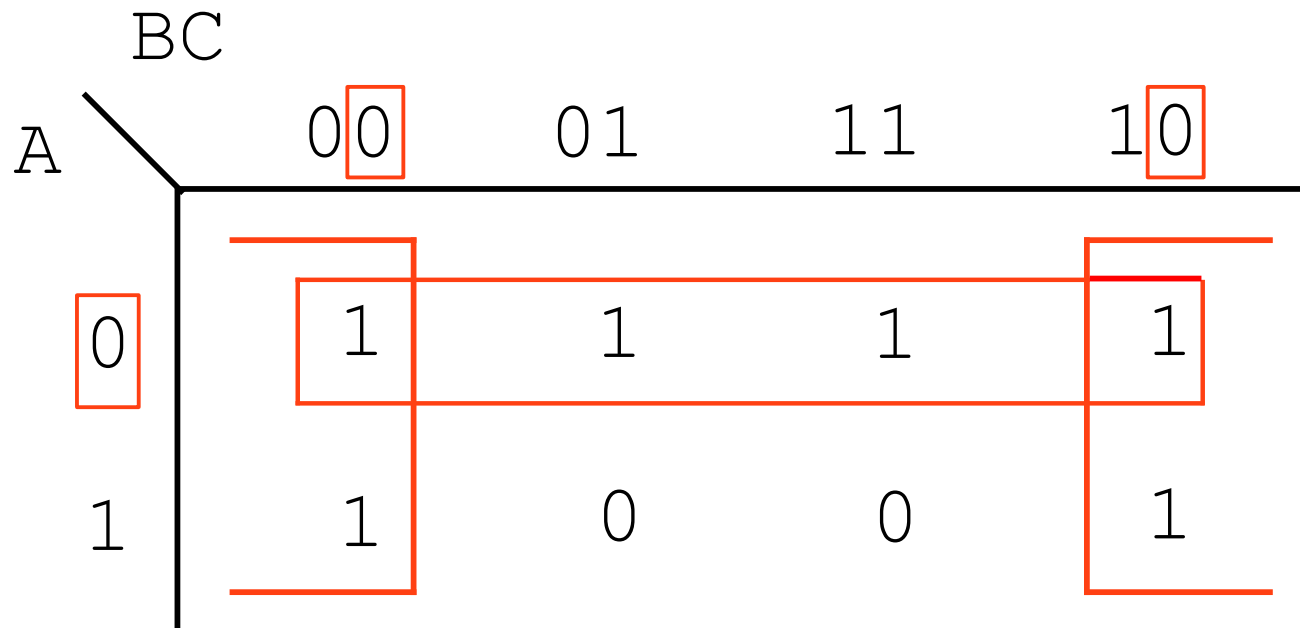
- Save the ones that stay the same in a group, discarding the rest
- This must be done for each group



$$R = \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}C + \bar{A}BC + \bar{A}B\bar{C} + A\bar{B}\bar{C} + AB\bar{C}$$

A	B	C	R
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	0

$$R = \bar{A} + \bar{C}$$



- Save the ones that stay the same in a group, discarding the rest
- This must be done for each group

# Four Variable Example

$$R = !A!B!C!D + !A!B!CD + !A!BC!D + !ABC!D + A!B!C!D + A!B!CD + A!BC!D$$

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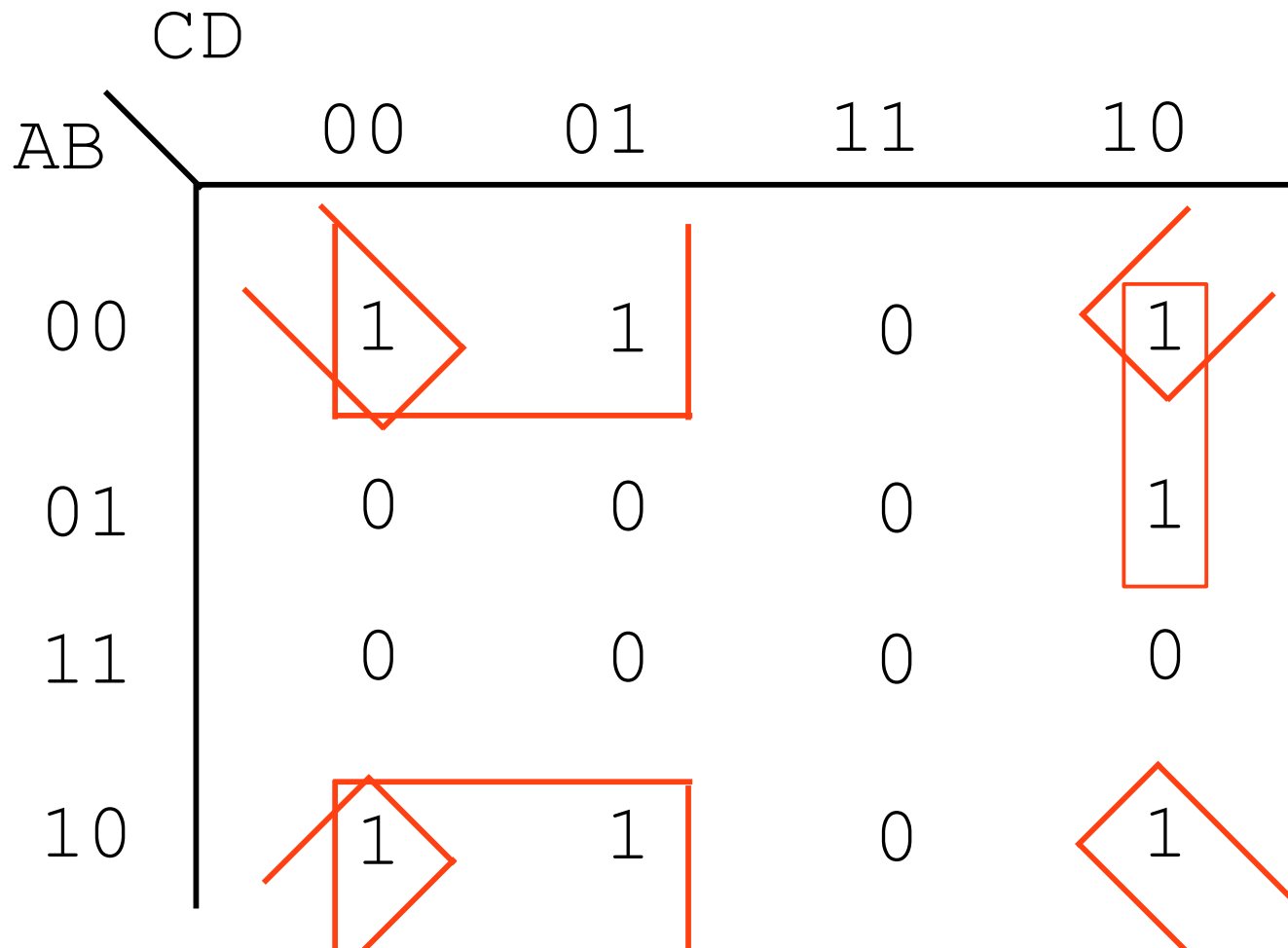
-Take this formula

$$R = !A!B!C!D + !A!B!CD + !A!BC!D + !ABC!D + A!B!C!D + A!B!CD + A!BC!D$$

		CD			
		00	01	11	10
AB	00	1	1	0	1
	01	0	0	0	1
	11	0	0	0	0
	10	1	1	0	1

-For space reasons, we go directly to the K-map

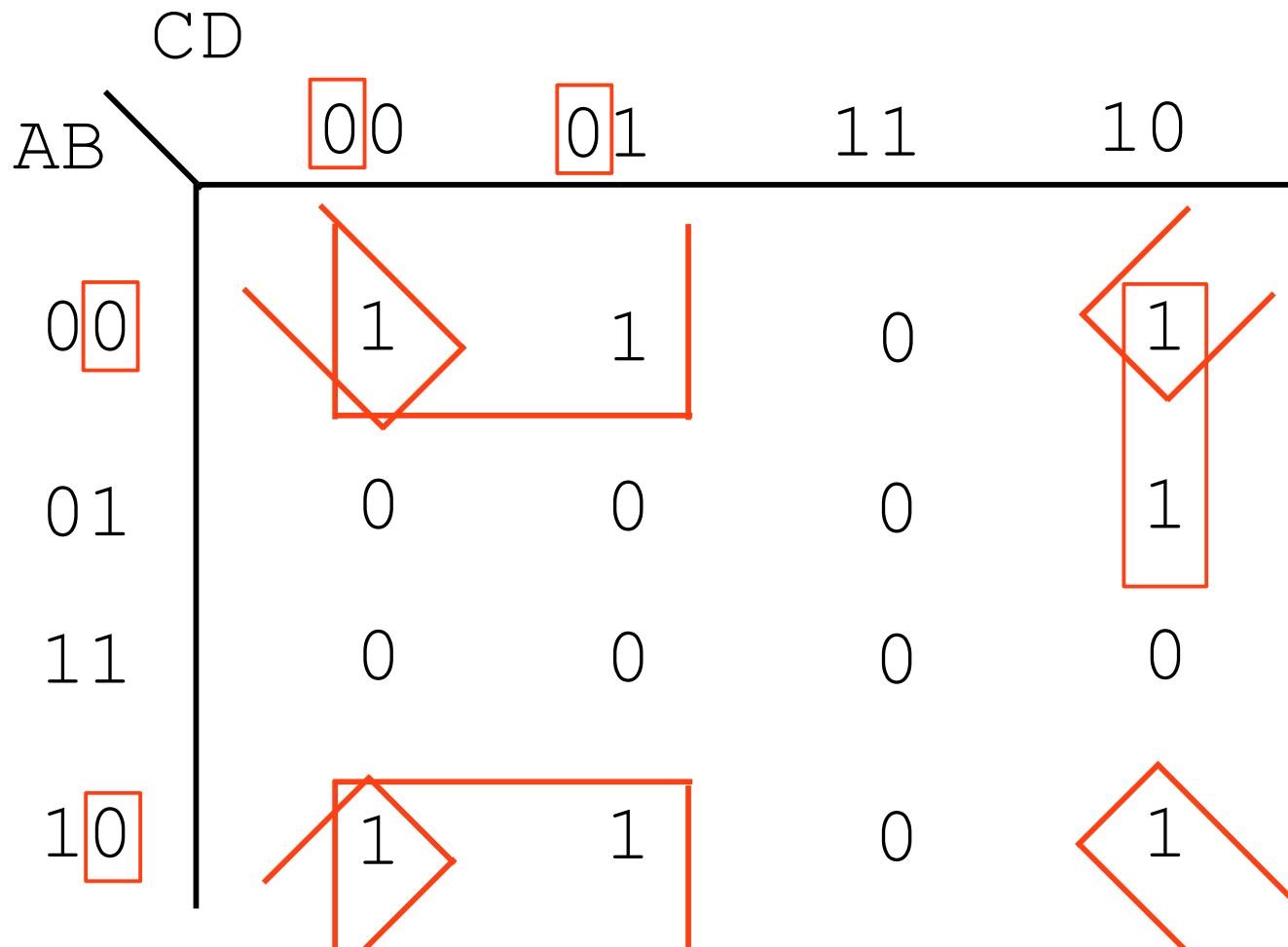
$$R = !A!B!C!D + !A!B!CD + !A!BC!D + !ABC!D + A!B!C!D + A!B!CD + A!BC!D$$



- Group things up
- The edges logically wrap around!
- Groups may overlap each other

$$R = !A!B!C!D + !A!B!CD + !A!BC!D + !ABC!D + A!B!C!D + A!B!CD + A!BC!D$$

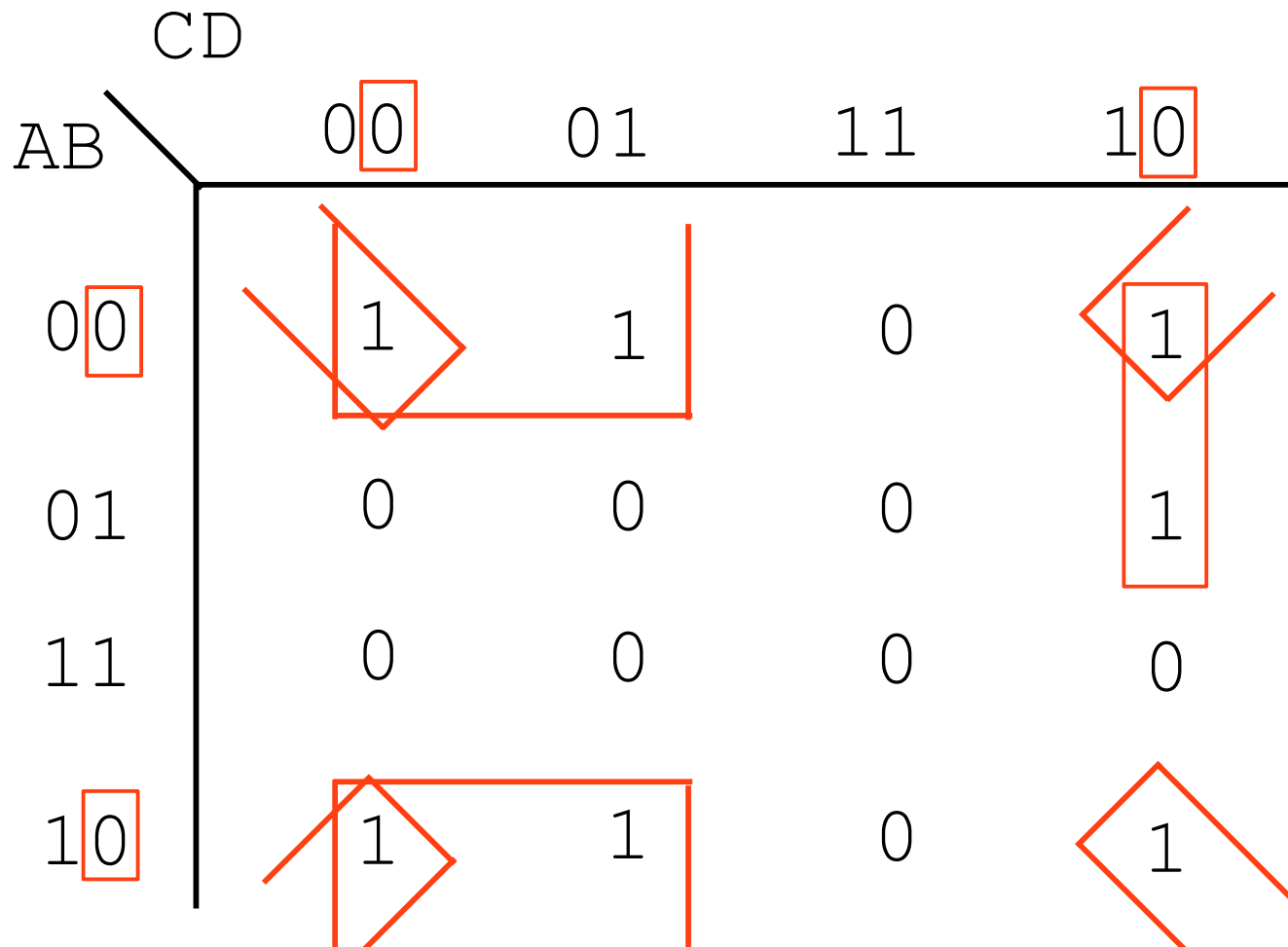
$$R = !B!C$$



- Look at the bits that don't change
- First for the cube

$$R = !A!B!C!D + !A!B!CD + !A!BC!D + !ABC!D + A!B!C!D + A!B!CD + A!BC!D$$

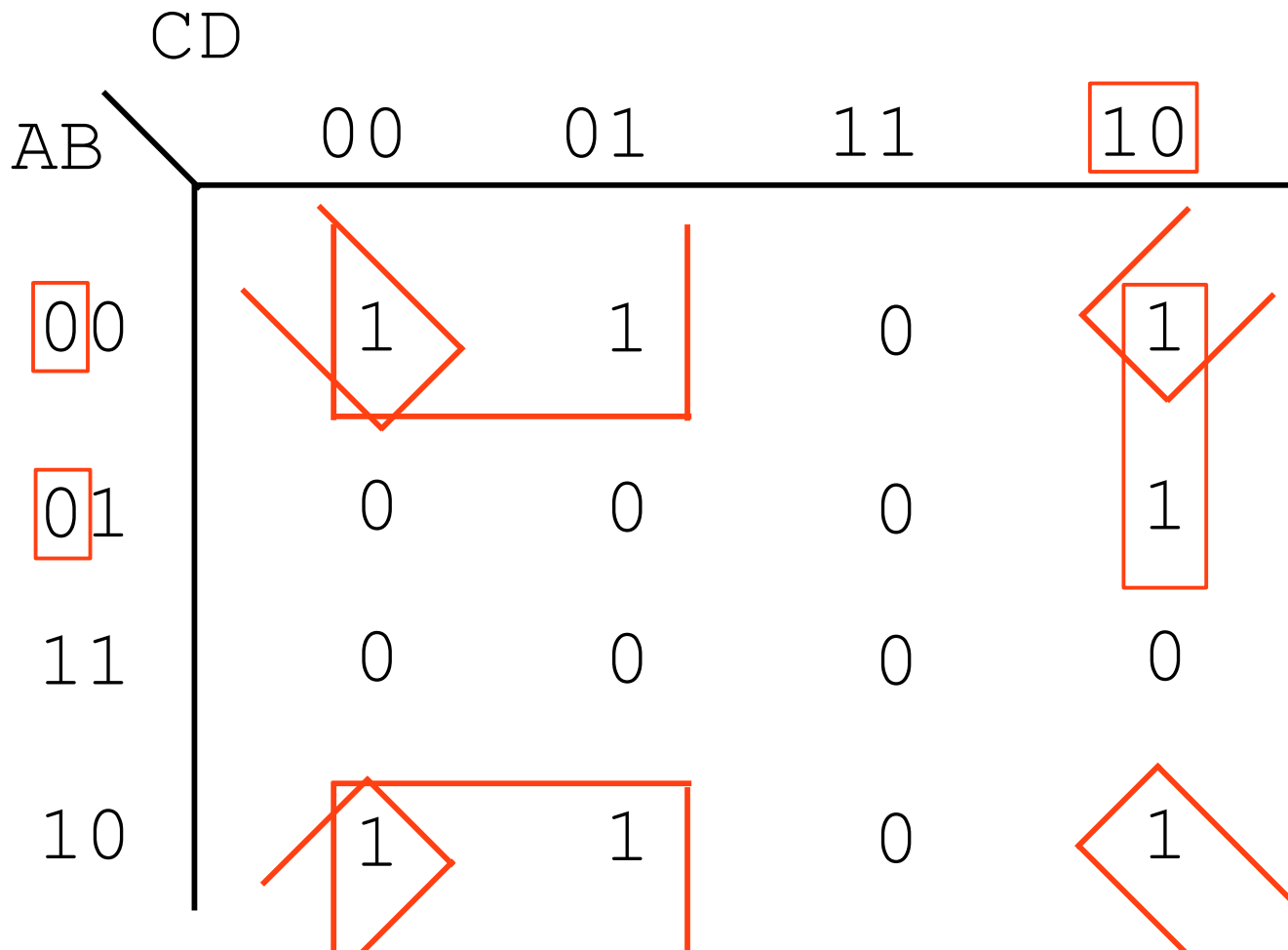
$$R = !B!C + !B!D$$



- Look at the bits that don't change
- Second for the cube on the edges

$$R = !A!B!C!D + !A!B!CD + !A!BC!D + !ABC!D + A!B!C!D + A!B!CD + A!BC!D$$

$$R = !B!C + !B!D + !AC!D$$



- Look at the bits that don't change
- Third for the line



# K-Map Rules in Summary (I)

- Groups can contain only 1s
- Only 1s in adjacent groups are allowed (no diagonals)
- The number of 1s in a group must be a power of two (1,2,4,8...)
- The groups must be as large as legally possible

# K-Map Rules in Summary (2)

- All 1s must belong to a group, even if it's a group of one element
- Overlapping groups are permitted
- Wrapping around the map is permitted
- Use the fewest number of groups possible

# Revisiting Problem

$$!A!BC + A!B!C + !ABC + !AB!C + A!BC$$

# Revisiting Problem

$$R = \neg A \neg B C + A \neg B \neg C + \neg A B C + \neg A B \neg C + A \neg B C$$

A	B	C	R
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	0
1	1	1	0

# Revisiting Problem

$$R = \neg A \neg B C + A \neg B \neg C + \neg A B C + \neg A B \neg C + A \neg B C$$

A	B	C	R
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	0
1	1	1	0

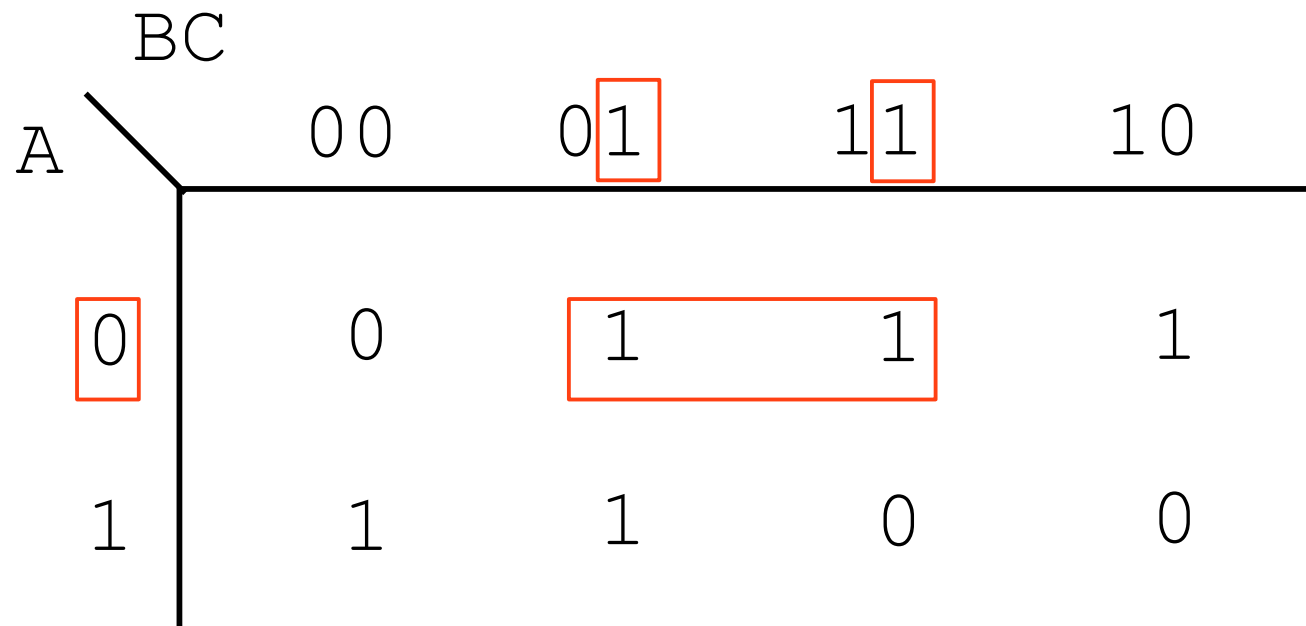
		BC			
A		00	01	11	10
	0	0	1	1	1
	1	1	1	0	0

# Revisiting Problem

$$R = \neg A \neg B C + A \neg B \neg C + \neg A B C + \neg A B \neg C + A \neg B C$$

A	B	C	R
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	0
1	1	1	0

$$R = \neg A C$$



# Revisiting Problem

$$R = \neg A \neg B C + A \neg B \neg C + \neg A B C + \neg A B \neg C + A \neg B C$$

A	B	C	R
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	0
1	1	1	0

$$R = \neg A C + A \neg B$$

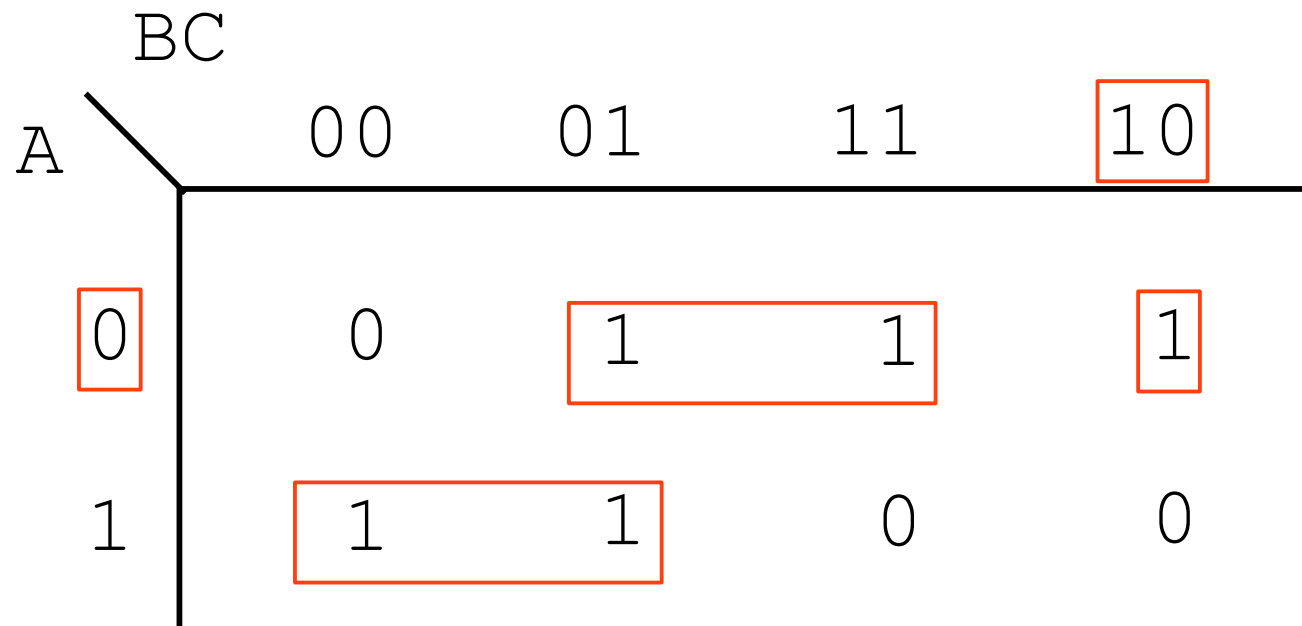
		BC			
A		00	01	11	10
0		0	1 1	1	1
1		1 1	1	0	0

# Revisiting Problem

$$R = \neg A \neg B C + A \neg B \neg C + \neg A B C + \neg A B \neg C + A \neg B C$$

A	B	C	R
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	0
1	1	1	0

$$R = \neg A C + A \neg B + \neg A B \neg C$$





# Difference

- Algebraic solution:  $\bar{B}C + A\bar{B}\bar{C} + \bar{A}B$
- K-map solution:  $\bar{A}C + A\bar{B} + \bar{A}B\bar{C}$
- Question: why might these differ?

# Difference

- Algebraic solution:  $\overline{B}C + A\overline{B}\overline{C} + \overline{A}B$
- K-map solution:  $\overline{A}C + A\overline{B} + \overline{A}B\overline{C}$
- Question: why might these differ?
  - Both are *minimal*, in that they have the fewest number of products possible
  - Can be multiple minimal solutions

# Difference

Algebraic solution:  $\neg BC + A\neg B\neg C + \neg AB$

K-map solution:  $\neg AC + A\neg B + \neg AB\neg C$

		BC			
		00	01	11	10
A	0	0	1	1	1
	1	1	1	0	0

-If we take our k-map from before with the grouping we chose, we get this particular solution

# Difference

Algebraic solution:  $\bar{B}C + A\bar{B}\bar{C} + \bar{A}B$

K-map solution:  $\bar{B}C + A\bar{B}\bar{C} + \bar{A}B$

		BC			
		00	01	11	10
A	0	0	1	1	1
	1	1	1	0	0

-If, however, we choose a different (also valid) grouping, we get the same solution as we did algebraically

# Don't Cares

Don't cares in a Karnaugh map, or truth table, may be either **1s** or **0s**, as long as we don't care what the output is for an input condition we never expect to see. We plot these cells with an asterisk, \*, among the normal **1s** and **0s**.

When forming groups of cells, treat the don't care cell as either a **1** or a **0**, or ignore the don't cares. This is helpful if it allows us to form a larger group than would otherwise be possible without the don't cares. There is no requirement to group all or any of the don't cares.

Only use them in a group if it simplifies the logic.

		BC			
A		00	01	11	10
0		0	0	0	X
1		0	1	X	X